

On a theorem of Lascar

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We denote by \mathbb{C} be a big saturated model (the monster model). If R, S are binary relations on \mathbb{C} the product RS is the class of all pairs (a, c) for which there is a $b \in \mathbb{C}$ such that aRb and bSc . The smallest type-definable relation which contains R is denoted by \overline{R} . The smallest invariant equivalence relation which has a bounded number of classes is E_L , the relation of having the same strong Lascar type. The smallest bounded type-definable equivalence relation, E_{KP} , was introduced by Kim and Pillay.

The following theorem was proved by Lascar using the Lascar galois group:

Theorem 1 (D. Lascar)

$$\overline{E_L} E_L = E_{KP}$$

I will give another of this theorem.

A formula $\theta(x, y)$ is called thick if there is no infinite sequence of (a_i) such that $\neg\theta(a_i, a_j)$ for $i < j$. We denote by $\Theta(x, y)$ the relation which is defined by the set of all thick formulas.

It is well known that E_L is the transitive closure of Θ :

$$E_L = \Theta \cup \Theta^2 \cup \Theta^3 \cup \dots$$

We will prove Theorem 1 in the following slightly stronger form.

Theorem 2

$$\overline{E_L} \Theta = E_{KP}$$

Lemma 3 (Open mapping) *Let A be a set of parameters, a an element and $\theta(x, y)$ a thick formula, possibly with parameters from A . Then there is an L_A -formula ϕ in $\text{tp}(a/A)$ such that every type $p \in S(A)$ which contains ϕ can be realized by an element b such that $\mathbb{C} \models \theta(a, b)$.*

PROOF: We can assume that A is empty. Otherwise we name the elements of A .

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Let D be the class of all conjugates of a and D_0 a finite subset of D such that D is contained in the definable class

$$B = \{b \in \mathbb{C} \mid \mathbb{C} \models \theta(a_0, b) \text{ for some } a_0 \in D_0\}.$$

The set of all types $p \in S(\emptyset)$ which can be realized by an element of $\mathbb{C} \setminus B$ is closed. Therefore the set Φ of all p with only realizations in B is open. Since $D \subset B$, Φ contains $\text{tp}(a)$. Thus we find a $\phi \in \text{tp}(a)$ such that every p which contains ϕ can only be realized by elements of B .

Fix a $p \in S(\emptyset)$ which contains ϕ . Choose a realization $b \in B$ and $a_0 \in D_0$ such that $\mathbb{C} \models \theta(a_0, b)$. Since a_0 has the same type as a we find a b' such that ab' has the same type as a_0b . Then b' realizes p and $\mathbb{C} \models \theta(a, b')$. This proves the Lemma.

Let R and S be two invariant relations on C . It is easy to see¹ that \overline{RS} is always contained in $\overline{R}\overline{S}$. The converse inclusion is not generally true: Take as a model a set with a sequence of named elements $0, 1, 2, \dots$. Take for R the set of all pairs $(0, 1), (0, 3), (0, 5), \dots$ and for S the set of all pairs $(2, 0), (4, 0), (6, 0), \dots$. Then \overline{RS} contains $(0, 0)$ and $\overline{R}\overline{S}$ is empty.

Of course, R and S are not connected in the following sense:

Definition 4 *Two invariant relation R and S are called connected if there is a complete type p over \emptyset such that both, $R(x, y)$ and $S(y, z)$, imply $p(y)$.*

EXAMPLE: Look at the group $G = \mathbb{R} \times \mathbb{R}$ with the lexicographical ordering. Forget everything except the ordering, addition with $(1, 0)$ and addition with $(0, 1)$. Define

$$R(x, y) \Leftrightarrow \bigvee_{n \in \omega} y < x + (0, n).$$

Since there is only one type over the empty set, R and R are connected. Also R is transitive, while

$$\overline{R}(x, y) \Leftrightarrow \bigwedge_{n \in \omega} y < x + (1, -n)$$

is not. Whence $\overline{R}\overline{R} \neq \overline{R} = \overline{RR}$.

Lemma 5 *Assume the invariant relations R and S to be connected². Then*

$$\overline{RS} \subset \overline{R}\overline{S}\Theta \tag{1}$$

$$\overline{RS} \subset \Theta\overline{RS} \tag{2}$$

$$\overline{R}\overline{S} \subset \Theta\overline{R}\overline{S}\Theta \tag{3}$$

¹Note that the product of two type-definable relations is again type-definable.

²For (1) (and similarly for (2)) we only need that the first components of all pairs in S realize the same type.

PROOF: We prove first (1). Assume $(\overline{RS})(a, c)$. Since $(\overline{RS}\Theta)(x, z')$ can be axiomatized

$$\{\exists z(\psi(x, z) \wedge \theta(z, z')) \mid \psi \in \overline{RS}, \theta \text{ thick}\}$$

we have to show that for all $\psi(x, z) \in \overline{RS}$ and all thick θ there is a c' such that $\overline{RS}(a, c')$ and $\theta(c', c)$. Let b be such that $R(a, b)$ and $\overline{S}(b, c)$. If we apply Lemma 3 to $\text{tp}(c/b)$ we obtain a formula $\phi(z, b)$ such that every type over b which realizes ϕ can be realized by an element c' which satisfies $\theta(c', c)$. Since $\overline{S}(b, c)$, and R and S are connected, there is a c' which realizes ϕ and satisfies $S(b, c')$. By the choice of ϕ we can choose c' in such a way that $\theta(c', c)$.

The proof of (2) is symmetrical.

(3) follows from (1) and (2):

$$\overline{RS} \subset \overline{RS}\Theta \subset \overline{\Theta RS}\Theta = \Theta \overline{RS}\Theta$$

PROOF OF THEOREM 2: Fix a complete type. **First we prove the theorem restricted to the type p .** We restrict the the meaning of E_L , \overline{E}_L , Θ and E_{KP} the the realization set of p . We can then apply the last lemma. Since $\overline{E}_L\Theta$ is type-definable it suffices to prove that $\overline{E}_L\Theta$ is transitive.

We have by the lemma $E_L\overline{E}_L \subset \overline{E}_L\Theta$ and $\overline{E}_L E_L \subset \Theta\overline{E}_L$. This gives

$$\overline{E}_L\Theta \subset \overline{E}_L E_L \subset \Theta\overline{E}_L \subset E_L\overline{E}_L \subset \overline{E}_L\Theta$$

and all four terms are equal. Part (3) of the Lemma gives

$$\overline{E}_L \overline{E}_L \subset \Theta\overline{E}_L\Theta = \overline{E}_L E_L.$$

Whence $\overline{E}_L\Theta$ is transitive:

$$\overline{E}_L\Theta\overline{E}_L\Theta = \overline{E}_L \overline{E}_L E_L = \overline{E}_L E_L^2 = \overline{E}_L E_L = \overline{E}_L\Theta.$$

Now the general case: Let for every complete type $\overline{E}_L(p)$ be the closure of $E_L \cap p^2$, $\Theta(p) = \Theta \cap p^2$ and $E_{KP}(p)$ be the finest bounded type-definable equivalence relation on p . We have proved that

$$\overline{E}_L(p)\Theta(p) = E_{KP}(p).$$

Since $E_{KP}(p) = E_{KP} \cap p^2$ this implies

$$\bigcup_p \overline{E}_L(p)\Theta = E_{KP}.$$

But $\bigcup_p \overline{E}_L(p) \subset \overline{E}_L$ and the theorem is proved.