## An exposition of the compactness of $L(Q^{cf})$

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#### Abstract

We give an exposition of the compactness of  $L(Q^{cf})$ , for any set C of regular cardinals.

### 1 Introduction

We present here a new and short exposition of the proof of the compactness of the logic  $L(Q_C^{\text{cf}})$ , first-order logic extended by the cofinality quantifier  $Q_C^{\text{cf}}$ , where C is a class of regular cardinals. The logic and the proof of compactness are due to S. Shelah. The Compactness Theorem was stated and proved in [7], but this article is not self-contained and some fundamental steps of the proof must be found in the earlier article [6]. The interested reader consulting these two articles will soon realise that the structure of the proof is not completely transparent and that to fully understand the details requires a lot of work.

The most popular case of the cofinality quantifier is the logic  $L(Q_{\omega}^{\text{cf}})$  of the quantifier of cofinality  $\omega$ , that is,  $C = \{\omega\}$ . Our motivation comes from the application of  $L(Q_{\omega}^{\text{cf}})$  in [1] to an old problem on expandability of models. An anonymous referee of a preliminary version of [1] did not accept the validity (in ZFC) of the compactness proof presented in [7], apparently confused by the assumption of the existence of a weakly compact cardinal made at the beginning of the article. The assumption only applies to a previous result on a logic stronger than first-order logic even for countable models.

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Our proof of compactness of  $L(Q_C^{cf})$  uses some ideas of [7], but it is more in the spirit of Keisler's proof in [4] of countable compactness of the logic  $L(Q_1)$  with the quantifier of uncountable cardinality. However we use a simpler notion of weak model. J. Väänänen in the last chapter of [8] offers also a proof of compactness of  $L(Q_{\omega}^{cf})$  in Keisler's style, but it is incomplete and only gives countable compactness (see I. Hodkinson's review in [3]).

There are some other proofs in the literature, but also unsatisfactory. The proof by H-D. Ebbinghaus in [2], based on a set-theoretical translation, is just an sketch and the proof of J.A. Makowsky and S. Shelah in [5] only replaces part of Shelah's argument in [7] by a different reasoning and does not include all details.

### 2 Connections

For a linear ordering (X, <) we use the expressions

$$\exists^{\mathrm{cf}} x A(x), \text{ and } \forall^{\mathrm{cf}} x A(x)$$

for  $\forall x' \exists x \ (x' \leq x \land A(x))$ , and  $\exists x' \forall x \ (x' \leq x \to A(x))$ , respectively.

**Definition.** Let X and Y be two linear orderings. A connection between X and Y is a relation  $G \subset X \times Y$  with satisfies

$$\exists^{\mathrm{cf}} x \,\forall^{\mathrm{cf}} y \, G(x, y) \text{ and} \tag{1}$$

$$\exists^{\mathrm{cf}} y \,\forall^{\mathrm{cf}} x \,\neg G(x, y). \tag{2}$$

Note that X and Y cannot be connected if X or Y has a last element.

# **Remark 2.1.** 1. If X has no last element, the relation $x \le y$ connects X with itself.

- 2. If G connects X and Y, then  $\neg G^{-1} = \{(y, x) \mid \neg G(x, y)\}$  connects Y and X.
- 3. If G connects X and Y, and H connects Y and Z, then

$$K = \left\{ (x, z) \mid \exists y' \left( \forall y \left( y' \le y \to G(x, y) \right) \land H(y', z) \right) \right\}$$

connects X and Z.

*Proof.* We will not make use of this remark, but we give a proof of 3, never-theless.

 $\exists^{cf} x \forall^{cf} z \ K(x,z)$ : If x' is given, there are x and y' that  $x' \leq x$  and  $y' \leq y \rightarrow G(x,y)$  for all y. If we choose y' large enough, there is also a z' such that  $z' \leq z \rightarrow H(y',z)$  for all z. This shows that  $z' \leq z \rightarrow K(x,z)$  for all z.

 $\exists^{cf} z \forall^{cf} x \neg K(x,z) : \text{ If } z' \text{ is given, we find } z, y' \text{ and } x' \text{ such that } z' \leq z \text{ and } and for all x and y we have <math>y' \leq y \rightarrow \neg H(y,z)$  and  $x' \leq x \rightarrow \neg G(x,y')$ . Now this implies that  $x' \leq x \rightarrow \neg K(x,z)$  for all x. To see this assume  $x' \leq x$ . We will show that  $\forall y \ (y'' \leq y \rightarrow G(x,y)) \land H(y'',z)$  is wrong for all y''. Indeed, if  $y'' \leq y'$ , this follows from  $\neg G(x,y')$ . And if  $y' \leq y''$ , we have  $\neg H(y'',z)$ .

**Remark 2.2.** If X and Y are connected by G, then also by

$$G' = \Big\{ (x,y) \ \Big| \ \exists x' \ \big( x \le x' \land \forall y' \ (y \le y' \to G(x',y')) \big) \Big\}.$$

G' is antitone in x and monotone in y.

*Proof.* It is easy to see that  $G^{\text{anti}} = \{(x, y) \mid \exists x' \ (x \leq x' \land G(x', y))\}$  connects X and Y and is antitone in x. Now set

$$G' = (\neg ((\neg G^{-1})^{\text{anti}})^{-1})^{\text{anti}}.$$

**Lemma 2.3.** Two linear orders without last element are connected if and only if they have the same cofinality.

*Proof.* If  $cf(X) = cf(Y) = \kappa$ , choose two increasing cofinal sequences  $(x_{\alpha} \mid \alpha < \kappa)$  and  $(y_{\alpha} \mid \alpha < \kappa)$  in X and Y. Then

$$G = \{(x, y) \mid \exists \alpha \ (x \le x_\alpha \land y_\alpha \le y)\}$$

connects X and Y.<sup>1</sup>

For the converse assume that  $cf(X) = \kappa$ , and that G connects X and Y. Choose a cofinal sequence  $(x_{\alpha} \mid \alpha < \kappa)$  in X and elements  $y_{\alpha}$  in Y such that  $y_{\alpha} \leq y \rightarrow G(x_{\alpha}, y)$  for all y. Then the  $y_{\alpha}$  are cofinal in Y. To see this let y be an element of Y. Since the  $x_{\alpha}$  are cofinal, we have  $\neg G(x_{\alpha}, y)$  for some  $\alpha$ . It follows that  $y < y_{\alpha}$ .

<sup>&</sup>lt;sup>1</sup>It suffices to assume that the  $y_{\alpha}$  are increasing. Also one can use  $G = \{(x_{\alpha}, y) \mid y_{\alpha} \leq y\}$ .

**Lemma 2.4.** Assume that  $G \subset X \times Y$  satisfies

$$\exists^{\mathrm{cf}} x \exists y \ G(x, y) \tag{3}$$

$$\forall y' \,\exists x' \,\forall xy \ (x' \le x \land y \le y') \to \neg G(x, y). \tag{4}$$

Then  $G' = \{(x, y) \mid \exists y' \ (y' \leq y \land G(x, y'))\}$  connects X and Y.

Note that a connecting G which is monotone in y satisfies (3) and (4).

*Proof.* This is a straightforward verification.

### 3 The Main Lemma

Consider a *L*-structure *M* with two (parametrically) definable linear orderings,  $<_{\varphi}$  and  $<_{\psi}$  of its universe, both without last element. We say that  $\varphi$  and  $\psi$  are *definably connected* if there is a definable connection between  $(M, <_{\varphi})$  and  $(M, <_{\psi})$ .

**Lemma 3.1.** If  $\varphi$  and  $\psi$  are not definably connected, and c is a new constant, the theory

 $T' = \operatorname{Th}(M, m)_{m \in M} \cup \{m <_{\varphi} c \mid m \in M\}$ 

does not isolate the partial type  $\Sigma(y) = \{n <_{\psi} y \mid n \in M\}.$ 

*Proof.* Assume that  $\gamma(c, y)$ , for some L(M)-formula  $\gamma(x, y)$ , isolates  $\Sigma(y)$  in T'. This means that

- 1.  $T' \cup \{\gamma(c, y)\}$  is consistent.
- 2.  $T' \vdash \gamma(c, y) \rightarrow n <_{\psi} y$  for all  $n \in M$ .

We show that the relation G defined by  $\gamma(x, y)$  has properties (3) and (4) of Lemma 2.4, where  $X = (M, <_{\varphi})$  and  $Y = (M, <_{\psi})$ . This will contradict the hypothesis of our Lemma.

That  $T' \cup \{\gamma(c, y)\}$  is consistent means that for all  $m \in M$  the theory  $\operatorname{Th}(M, m)_{m \in M}$  does not prove  $m \leq_{\varphi} c \to \neg \exists y \ \gamma(c, y)$ , which means that  $M \models \exists x (m \leq_{\psi} x \land \exists y \ \gamma(x, y))$ . This is exactly condition (3) of 2.4.

That  $T' \vdash \gamma(c, y) \to n <_{\psi} y$  means that there is an  $m \in M$  such that  $\operatorname{Th}(M, m)_{m \in M}$  proves  $(m \leq_{\varphi} c \land \gamma(c, y)) \to n <_{\psi} y$ , which means  $M \models$   $\forall xy \ (m \leq_{\varphi} x \land y \leq_{\psi} n \to \neg \gamma(x, y))$ . The existence of such m for all n is exactly condition (4) of 2.4. **Corollary 3.2.** Assume  $\kappa$  is regular,  $|M|, |L| \leq \kappa$ , and  $<_{\varphi}$  is a definable linear ordering of M without last element. Then there is an elementary extension N of M such that:

- 1. M is not  $<_{\varphi}$ -cofinal in N.
- 2. If  $<_{\psi}$  is a definable linear ordering of M of cofinality  $\kappa$ , and  $\psi$  and  $\varphi$  are not definably connected, then M is  $<_{\psi}$ -cofinal in N.

Proof. Let c be a new constant and let  $T' = \operatorname{Th}(M, m)_{m \in M} \cup \{m <_{\varphi} c \mid m \in M\}$ . By Lemma 3.1, T' does not isolate any of the types  $\Sigma_{\psi}(y) = \{n <_{\psi} y \mid n \in M\}$ . By the form of the types and regularity of  $\kappa$ , for any  $<_{\psi}$  of cofinality  $\kappa$  the type  $\Sigma_{\psi}(y)$  cannot be isolated neither by means of a set of  $< \kappa$  formulas. By the  $\kappa$ -Omitting Types Theorem, there is a model of T' omitting all types  $\Sigma_{\psi}(y)$  for any  $<_{\psi}$  of cofinality  $\kappa$ . This gives the elementary extension N.

This corollary applies in particular to the case  $\kappa = \omega$ . Here the assumption on the cofinality of  $\langle \psi \rangle$  is not needed since it is the only possible cofinality in a countable model, and the Omitting Types Theorem used in the proof is the ordinary one for countable languages and countably many non-isolated types.

### 4 Completeness

For a language L let  $L(Q^{cf})$  be the set of formulas which are built like firstorder formulas but using an additional two-place quantifier  $Q^{cf}xy \varphi$ , for different variables x and y. Let C be class a of regular cardinals and M an L-structure. For a binary relation R on M, we write "cf  $R \in C$ " for "R is a linear ordering of M, without last element and cofinality in C".

The satisfaction relation  $\models_C$  for *L*-structures M,  $L(Q^{cf})$ -formulas  $\psi(\bar{z})$ , and tuples  $\bar{c}$  of elements of M is defined inductively, where the  $Q^{cf}$ -step is

$$M \models_C Q^{\mathrm{cf}} xy \,\varphi(x, y, \bar{c}) \, \Leftrightarrow \, \mathrm{cf} \left\{ (a, b) \mid M \models_C \varphi(a, b, \bar{c}) \right\} \in C.$$

We say that M is a C-model of T, a set of  $L(Q^{cf})$ -sentences, if  $M \models_C \psi$  for all  $\psi \in T$ .

A weak structure  $M^* = (M, ...)$  is an  $L^*$ -structure, where  $L^*$  is an extension of L by an *n*-ary relation  $R_{\varphi}$  for every  $L(Q^{\text{cf}})$ -formula  $\varphi(x, y, z_1, ..., z_n)$ . Satisfaction is defined using the rule

$$M^* \models Q^{\mathrm{cf}} xy \ \varphi(x, y, \bar{c}) \ \Leftrightarrow \ M^* \models R_{\varphi}(\bar{c}).$$

In weak structures every  $L(Q^{cf})$ -formula is equivalent to a first-order  $L^*$ -formula, and conversely. So the  $L(Q^{cf})$ -model theory of weak structures is the same as their first-order model theory.

Note that the C-semantics of M is given by the semantics of the weak structure  $M^*$  if one sets

$$M^* \models R_{\varphi}(\bar{c}) \iff M \models_C Q^{\mathrm{cf}} xy \ \varphi(x, y, \bar{c}).$$

The following lemma is clear:

**Lemma 4.1.** A weak structure  $M^*$  describes the C-semantics of M if and only if

$$M^* \models Q^{\mathrm{cf}} xy \ \varphi(x, y, \bar{c}) \ \Leftrightarrow \ \mathrm{cf} \left\{ (a, b) \mid M^* \models \varphi(a, b, \bar{c}) \right\} \in C$$

for all  $\varphi$  and  $\bar{c}$ .

The following property of weak structures  $M^*$  can be expressed by a set SA of  $L(Q^{cf})$  sentences (the Shelah Axioms):

If the  $L(Q^{cf})(M)$ -formula  $\varphi(x, y)$  satisfies  $M^* \models Q^{cf}xy \ \varphi(x, y)$  then  $\varphi$  defines a linear ordering  $<_{\varphi}$  without last element. Furthermore, if  $\psi(x, y)$  defines a linear ordering  $<_{\psi}$  and  $M^* \models \neg Q^{cf}xy \ \psi(x, y)$ , there is no definable connection between  $(M, <_{\varphi})$  and  $(M, <_{\psi})$ .

Lemma 4.2. L-structures with the C-semantics are models of SA.

*Proof.* This follows from Lemma 2.3.

**Theorem 4.3.** Let C be a non-empty class of regular cardinals, different from the class of all regular cardinals. An  $L(Q^{cf})$ -theory T has a C-model if and only if  $T \cup SA$  has a weak model.

*Proof.* One direction follows from Lemma 4.2. For the other direction assume that  $T \cup SA$  has a weak model.

Claim 1: If L is countable, T has a  $\{\omega\}$ -model of cardinality  $\omega_1$ .

Proof. Let  $M_0^*$  be countable weak model of  $T \cup SA$ . Consider a linear ordering  $<_{\varphi}$  without last element and  $M_0^* \models \neg Q^{\mathrm{cf}} xy \varphi$ . Then by Corollary 3.2 for  $\kappa = \omega$  and the axioms SA, there is an elementary extension  $M_1^*$  such that  $M_0$  is not  $<_{\varphi}$ -cofinal in  $M_1$ , but  $<_{\psi}$ -cofinal in  $M_1$  for every  $\psi$  with  $M_0^* \models Q^{\mathrm{cf}} xy \psi$ . We may assume that  $M_1^*$  is countable. Continuing in this manner, taking unions at limit stages, one constructs an elementary chain of countable weak models  $M_0^* \prec M_1^* \cdots$  of length  $\omega_1$  with union  $M^*$ , such that

- 1. If  $<_{\varphi}$  is a linear ordering of  $M^*$  without last element and  $M^* \models \neg Q^{\text{cf}} xy \varphi$ , and if the parameters of  $\varphi$  are in  $M_{\alpha}$ , then for uncountably many  $\beta \ge \alpha$ ,  $M_{\beta}$  is not  $<_{\varphi}$ -cofinal in  $M_{\beta+1}$ .
- 2. If  $M^* \models Q^{cf} xy \psi$ , and the parameters of  $\varphi$  are in  $M_{\alpha}$ , then  $M_{\alpha}$  is  $<_{\psi}$ -cofinal in M.

It follows that, if  $M^* \models \neg Q^{cf} xy \varphi$ , then either  $\varphi$  does not define a linear ordering without last element, or  $<_{\varphi}$  has cofinality  $\omega_1$ . And, if  $M^* \models Q^{cf} xy \psi$ , then  $<_{\psi}$  has cofinality  $\omega$ . By Lemma 4.1 M is an  $\{\omega\}$ -model of the  $L(Q^{cf})$ theory of  $M^*$ , and whence an  $\{\omega\}$ -model of T. This proves Claim 1.

Let L' be the extension of L which has for every  $L(Q^{\text{cf}})$ -formula  $\varphi(x, y, \bar{z})$ a new relation symbol  $V_{\varphi}$  of arity  $2 + 2 \cdot |\bar{z}|$ . Let SK be the set of axioms which state that if  $\varphi(x, y, \bar{c}_1)$  and  $\varphi(x, y, \bar{c}_2)$  define linear orderings without last elements, and

$$Q^{\mathrm{cf}}xy\,\varphi(x,y,\bar{c}_1) \leftrightarrow Q^{\mathrm{cf}}xy\,\varphi(x,y,\bar{c}_2),$$

then  $V_{\varphi}(x, y, \bar{c}_1, \bar{c}_2)$  defines a connection between the two orderings.

Claim 2:  $T \cup SA \cup SK$  has a weak model.

Proof: By compactness we may assume that L is countable. Then T has an  $\{\omega\}$ -model M of cardinality  $\omega_1$ , by Claim 1. If  $\varphi(x, y, \bar{c}_1)$  and  $\varphi(x, y, \bar{c}_2)$ define linear orderings without last element, and  $M \models_C Q^{cf} xy \varphi(x, y, \bar{c}_1) \leftrightarrow$  $Q^{cf} xy \varphi(x, y, \bar{c}_2)$ , then the two orderings have the same cofinality, namely  $\omega$ or  $\omega_1$ , and there is a connection between them by Lemma 2.3. This proves Claim 2.

To prove the theorem, we choose two regular cardinals  $\lambda, \kappa$  such that  $|L| \leq \kappa$  and either  $\lambda \notin C$  and  $\kappa \in C$  or conversely. Let  $M_0^*$  be a weak model

of  $T \cup SA \cup SK$ . It  $M_0^*$  is finite, it is a *C*-model of *T* for trivial reasons<sup>2</sup>. Otherwise we may assume that  $M_0^*$  has cardinality  $\kappa$  and all  $L(Q^{cf})$ -definable linear orderings without last element have cofinality  $\kappa$ . Let us first assume that  $\lambda \notin C$  and  $\kappa \in C$ .

Consider an  $L(Q^{cf})$ -definable linear ordering  $<_{\varphi}$  without last element and  $M_0^* \models \neg Q^{cf}xy \varphi$ . Then by Corollary 3.2 and the axioms SA, there is an elementary extension  $M_1^*$  such that  $M_0$  is not  $<_{\varphi}$ -cofinal in  $M_1$ , but  $<_{\psi}$ -cofinal in  $M_1$  for every  $L(Q^{cf})$ -formula  $\psi$  with  $M_0^* \models Q^{cf}xy\psi$ . We may assume that  $M_1^*$  has cardinal  $\kappa$ . The axioms SK imply that in  $M_1^*$  every  $L(Q^{cf})$ -definable linear ordering without last element is connected to a linear ordering defined in  $M_0$ , and so has also cofinality  $\kappa$ .

Continuing in this manner, taking unions at limit stages, one constructs an elementary chain of weak models  $M_0^* \prec M_1^* \cdots$  of length  $\lambda$  with union  $M^*$ , such that

- 1. If  $<_{\varphi}$  is an  $L(Q^{cf})$ -definable linear ordering  $<_{\varphi}$  of  $M^*$  without last element and  $M^* \models \neg Q^{cf} xy \varphi$ , and if the parameters of  $\varphi$  are in  $M_{\alpha}$ , then for  $\lambda$ -many  $\beta \ge \alpha$ ,  $M_{\beta}$  is not  $<_{\varphi}$ -cofinal in  $M_{\beta+1}$ .
- 2. If  $M^* \models Q^{cf}xy\psi$ , and the parameters of  $\varphi$  are in  $M_{\alpha}$ , then  $M_{\alpha}$  is  $<_{\psi}$ -cofinal in M.

It follows that, if  $M^* \models \neg Q^{cf} xy \varphi$ , then either  $\varphi$  does not define a linear ordering without last element, or  $<_{\varphi}$  has cofinality  $\lambda$ . And, if  $M^* \models Q^{cf} xy \psi$ , then  $<_{\psi}$  has cofinality  $\kappa$ . By Lemma 4.1  $M \upharpoonright L$  is an *C*-model of the  $L(Q^{cf})$ -theory of  $M^*$ , and whence a *C*-model of *T*.

The proof in the case  $\lambda \in C$  and  $\kappa \notin C$  is, mutatis mutandis, the same.

**Corollary 4.4.** For every class C of regular cardinals, the logic  $L(Q_C^{cf})$  is compact.

We have always assumed that whenever  $Q^{cf}xy\varphi(x, y, \bar{c})$ , the definable ordering  $<_{\varphi}$  linearly orders the universe. This is not exactly the assumption of Shelah in [7]: with his definition  $<_{\varphi}$  linearly orders  $\{x \mid \exists y \varphi(x, y, \bar{c})\}$ , the domain of  $\varphi$ . The results presented here, in particular completeness and compactness, also apply to this modification of the semantics, it suffices to add, for each such  $\varphi$ , new relation symbols  $R_{\varphi}$  and  $H_{\varphi}$ , and declare

<sup>&</sup>lt;sup>2</sup>SA is used here.

that for every  $\bar{c}$ ,  $R_{\varphi}(x, y, \bar{c})$  defines a linear ordering  $<'_{\varphi}$  on the universe and  $H_{\varphi}(x, y, \bar{c})$  connects  $<_{\varphi}$  and  $<'_{\varphi}$ . This gives compactness. For the formulation of completeness (Theorem 4.3) one must adapt the axioms SA to the new situation.

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