

# On a theorem of Lascar

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October 5, 2000

We denote by  $\mathbb{C}$  be a big saturated model (the monster model). If  $R, S$  are binary relations on  $\mathbb{C}$  the product  $RS$  is the class of all pairs  $(a, c)$  for which there is a  $b \in \mathbb{C}$  such that  $aRb$  and  $bSc$ . The smallest type-definable relation which contains  $R$  is denoted by  $\overline{R}$ . The smallest invariant equivalence relation which has a bounded number of classes is  $E_L$ , the relation of having the same strong Lascar type. The smallest bounded type-definable equivalence relation,  $E_{KP}$ , was introduced by Kim and Pillay.

The following theorem was proved by Lascar using the Lascar galois group:

**Theorem 1 (D. Lascar)**

$$\overline{E_L} E_L = E_{KP}$$

I will give another proof of this theorem.

A formula  $\theta(x, y)$  is called thick if there is no infinite sequence of  $(a_i)$  such that  $\neg\theta(a_i, a_j)$  for  $i < j$ . We denote by  $\Theta(x, y)$  the relation which is defined by the set of all thick formulas.

It is well known that  $E_L$  is the transitive closure of  $\Theta$ :

$$E_L = \Theta \cup \Theta^2 \cup \Theta^3 \cup \dots$$

We will prove Theorem 1 in the following slightly stronger form.

**Theorem 2**

$$\overline{E_L} \Theta = E_{KP}$$

**Lemma 3 (Open mapping)** *Let  $A$  be a set of parameters,  $a$  an element and  $\theta(x, y)$  a thick formula, possibly with parameters from  $A$ . Then there is an  $L_A$ -formula  $\phi$  in  $\text{tp}(a/A)$  such that every type  $p \in S(A)$  which contains  $\phi$  can be realized by an element  $b$  such that  $\mathbb{C} \models \theta(a, b)$ .*

PROOF: We can assume that  $A$  is empty. Otherwise we name the elements of  $A$ .

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\*I thank H. Casanovas and F. Wagner for a helpful discussion.

Let  $D$  be the class of all conjugates of  $a$  and  $D_0$  a finite subset of  $D$  such that  $D$  is contained in the definable class

$$B = \{b \in \mathbb{C} \mid \mathbb{C} \models \theta(a_0, b) \text{ for some } a_0 \in D_0\}.$$

The set of all types  $p \in S(\emptyset)$  which can be realized by an element of  $\mathbb{C} \setminus B$  is closed. Therefore the set  $\Phi$  of all  $p$  with only realizations in  $B$  is open. Since  $D \subset B$ ,  $\Phi$  contains  $\text{tp}(a)$ . Thus we find a  $\phi \in \text{tp}(a)$  such that every  $p$  which contains  $\phi$  can only be realized by elements of  $B$ .

Fix a  $p \in S(\emptyset)$  which contains  $\phi$ . Choose a realization  $b \in B$  and  $a_0 \in D_0$  such that  $\mathbb{C} \models \theta(a_0, b)$ . Since  $a_0$  has the same type as  $a$  we find a  $b'$  such that  $ab'$  has the same type as  $a_0b$ . Then  $b'$  realizes  $p$  and  $\mathbb{C} \models \theta(a, b')$ . This proves the Lemma.

Let  $R$  and  $S$  be two invariant relations on  $C$ . It is easy to see<sup>1</sup> that  $\overline{RS}$  is always contained in  $\overline{R}\overline{S}$ . The converse inclusion is not generally true: Take as a model a set with a sequence of named elements  $0, 1, 2, \dots$ . Take for  $R$  the set of all pairs  $(0, 1), (0, 3), (0, 5), \dots$  and for  $S$  the set of all pairs  $(2, 0), (4, 0), (6, 0), \dots$ . Then  $\overline{RS}$  contains  $(0, 0)$  and  $\overline{R}\overline{S}$  is empty.

Of course,  $R$  and  $S$  are not connected in the following sense:

**Definition 4** *Two invariant relation  $R$  and  $S$  are called connected if there is a complete type  $p$  over  $\emptyset$  such that both,  $R(x, y)$  and  $S(y, z)$ , imply  $p(y)$ .*

EXAMPLE: Look at the group  $G = \mathbb{R} \times \mathbb{R}$  with the lexicographical ordering. Forget everything except the ordering, addition with  $(1, 0)$  and addition with  $(0, 1)$ . Define

$$R(x, y) \Leftrightarrow \bigvee_{n \in \omega} y < x + (0, n).$$

Since there is only one type over the empty set,  $R$  and  $R$  are connected. Also  $R$  is transitive, while

$$\overline{R}(x, y) \Leftrightarrow \bigwedge_{n \in \omega} y < x + (1, -n)$$

is not. Whence  $\overline{R}\overline{R} \neq \overline{R} = \overline{RR}$ .

**Lemma 5** *Assume the invariant relations  $R$  and  $S$  to be connected<sup>2</sup>. Then*

$$\overline{RS} \subset \overline{R}\overline{S}\Theta \tag{1}$$

$$\overline{RS} \subset \Theta\overline{RS} \tag{2}$$

$$\overline{R}\overline{S} \subset \Theta\overline{R}\overline{S}\Theta \tag{3}$$

<sup>1</sup>Note that the product of two type-definable relations is again type-definable.

<sup>2</sup>For (1) (and similarly for (2)) we only need that the first components of all pairs in  $S$  realize the same type.

PROOF: We prove first (1). Assume  $(\overline{RS})(a, c)$ . Since  $(\overline{RS}\Theta)(x, z')$  can be axiomatized

$$\{\exists z(\psi(x, z) \wedge \theta(z, z')) \mid \psi \in \overline{RS}, \theta \text{ thick}\}$$

we have to show that for all  $\psi(x, z) \in \overline{RS}$  and all thick  $\theta$  there is a  $c'$  such that  $\overline{RS}(a, c')$  and  $\theta(c', c)$ . Let  $b$  be such that  $R(a, b)$  and  $\overline{S}(b, c)$ . If we apply Lemma 3 to  $\text{tp}(c/b)$  we obtain a formula  $\phi(z, b)$  such that every type over  $b$  which realizes  $\phi$  can be realized by an element  $c'$  which satisfies  $\theta(c', c)$ . Since  $\overline{S}(b, c)$ , and  $R$  and  $S$  are connected, there is a  $c'$  which realizes  $\phi$  and satisfies  $S(b, c')$ . By the choice of  $\phi$  we can choose  $c'$  in such a way that  $\theta(c', c)$ .

The proof of (2) is symmetrical.

(3) follows from (1) and (2):

$$\overline{RS} \subset \overline{RS}\Theta \subset \overline{\Theta RS}\Theta = \Theta \overline{RS}\Theta$$

PROOF OF THEOREM 2: Fix a complete type. **First we prove the theorem restricted to the type  $p$ .** We restrict the the meaning of  $E_L$ ,  $\overline{E}_L$ ,  $\Theta$  and  $E_{KP}$  the the realization set of  $p$ . We can then apply the last lemma. Since  $\overline{E}_L\Theta$  is type-definable it suffices to prove that  $\overline{E}_L\Theta$  is transitive.

We have by the lemma  $E_L\overline{E}_L \subset \overline{E}_L\Theta$  and  $\overline{E}_L E_L \subset \Theta\overline{E}_L$ . This gives

$$\overline{E}_L\Theta \subset \overline{E}_L E_L \subset \Theta\overline{E}_L \subset E_L\overline{E}_L \subset \overline{E}_L\Theta$$

and all four terms are equal. Part (3) of the Lemma gives

$$\overline{E}_L \overline{E}_L \subset \Theta\overline{E}_L\Theta = \overline{E}_L E_L.$$

Whence  $\overline{E}_L\Theta$  is transitive:

$$\overline{E}_L\Theta\overline{E}_L\Theta = \overline{E}_L \overline{E}_L E_L = \overline{E}_L E_L^2 = \overline{E}_L E_L = \overline{E}_L\Theta.$$

Now the general case: Let for every complete type  $\overline{E}_L(p)$  be the closure of  $E_L \cap p^2$ ,  $\Theta(p) = \Theta \cap p^2$  and  $E_{KP}(p)$  be the finest bounded type-definable equivalence relation on  $p$ . We have proved that

$$\overline{E}_L(p)\Theta(p) = E_{KP}(p).$$

Since  $E_{KP}(p) = E_{KP} \cap p^2$  this implies

$$\bigcup_p \overline{E}_L(p)\Theta = E_{KP}.$$

But  $\bigcup_p \overline{E}_L(p) \subset \overline{E}_L$  and the theorem is proved.